

Chapter 1

Library eo

```
Require Import Arith.
Require Import Ascii.
Require Import Bool.
Require Import Coq.Strings.Byte.
Import IFNOTATIONS.

Inductive string : Set :=
  | EmptyString : string
  | String : ascii → string → string.

Inductive EOExpr : Type :=
  | EUnit : EOExpr
  | EInt : nat → EOExpr
  | EBool : bool → EOExpr
  | EVar : string → EOExpr
  | EAssign : EOExpr → EOExpr → EOExpr
  | EPlus : EOExpr → EOExpr → EOExpr
  | EMinus : EOExpr → EOExpr → EOExpr
  | EMult : EOExpr → EOExpr → EOExpr
  | EDiv : EOExpr → EOExpr → EOExpr
  | EIf : EOExpr → EOExpr → EOExpr → EOExpr
  | ELambda : string → EOExpr → EOExpr
  | EClosure : string → EOExpr → EOEnv → EOExpr
  | EApply : EOExpr → EOExpr → EOExpr

with EOEnv : Type :=
  | EEmptyEnv : EOEnv
  | EBind : string → EOExpr → EOEnv → EOEnv.
```

Chapter 2

Library main

Require Import Bool.

Inductive **bool** :=

| true

| false.

Definition negb ($b : \mathbf{bool}$) : **bool** :=

 match b with

 | true \Rightarrow false

 | false \Rightarrow true

 end.

Theorem negb_negb: $\forall (b : \mathbf{bool}), \text{negb} (\text{negb } b) = b$.

Proof.

 intros b .

 destruct b .

 + simpl. reflexivity.

 + simpl. reflexivity.

Qed.

Theorem negb_trice : $\forall (b : \mathbf{bool}), \text{negb} (\text{negb} (\text{negb } b)) = \text{negb } b$.

Proof.

 intros b .

 destruct b .

 + simpl. reflexivity.

 + simpl. reflexivity.

Qed.

Definition andb ($a b : \mathbf{bool}$) : **bool** :=

 match a, b with

 | true, true \Rightarrow true

 | -, - \Rightarrow false

 end.

Theorem `and_true_both_arg_true` : $\forall (a\ b : \mathbf{bool}),$
 $a = \mathbf{true} \rightarrow b = \mathbf{true} \rightarrow \mathbf{andb}\ a\ b = \mathbf{true}.$

Proof.

```
intros a b Ha Hb.  
rewrite Ha.  
rewrite Hb.  
simpl. reflexivity.
```

Qed.

Theorem `and_true_otherside` : $\forall (a\ b : \mathbf{bool}),$
 $\mathbf{andb}\ a\ b = \mathbf{true} \rightarrow a = \mathbf{true} \wedge b = \mathbf{true}.$

Proof.

```
intros a b Ha.  
destruct a.  
destruct b.  
split. reflexivity.  
reflexivity.  
split. reflexivity.  
simpl in Ha.  
inversion Ha.  
destruct b.  
simpl in Ha.  
inversion Ha.  
simpl in Ha.  
inversion Ha.
```

Qed.

Theorem `andb_associative` : $\forall (a\ b\ c : \mathbf{bool}),$
 $\mathbf{andb}\ a\ (\mathbf{andb}\ b\ c) = \mathbf{andb}\ (\mathbf{andb}\ a\ b)\ c.$

Proof.

```
intros [] [] [];  
reflexivity.
```

Qed.

Theorem `andb_comutative` : $\forall a\ b, \mathbf{andb}\ a\ b = \mathbf{andb}\ b\ a.$

Proof.

```
intros [] []; simpl; reflexivity.
```

Qed.

Definition `orb`($a\ b : \mathbf{bool}$) : $\mathbf{bool} :=$

```
match a, b with  
| false, false  $\Rightarrow$  false  
| -, -  $\Rightarrow$  true  
end.
```

Theorem `andb_negb_orb` : $\forall (a\ b : \mathbf{bool}),$

$\text{negb (andb } a \text{ } b) = \text{orb (negb } a) \text{ (negb } b)$.

Proof.

```
intros [] []; simpl; reflexivity.
```

Qed.

Definition material($a \ b : \text{bool}$) : **bool** :=

```
match a,b with
```

```
| false, true  $\Rightarrow$  false
```

```
| -, -  $\Rightarrow$  true
```

```
end.
```

Theorem positive_material : $\forall a \ b : \text{bool}$,

$a = \text{true} \rightarrow \text{material } a \ b = \text{true}$.

Proof.

```
intros a b Ha.
```

```
destruct a eqn:E.
```

```
- simpl. destruct b; reflexivity.
```

```
- discriminate Ha.
```

```
Show Proof.
```

Qed.

Chapter 3

Library phi

Require Import **Relations.Relation_Definitions**.

Inductive **Process** :=
| Nil : **Process**
| Input : **nat** → (**nat** → **Process**) → **Process**
| Output : **nat** → **Process** → **Process**
| Sum : **Process** → **Process** → **Process**
| Res : **nat** → **Process** → **Process**.

Definition small_step : **relation Process** :=

```
fun p1 p2 =>  
  match p1 with  
  | Input x f =>  
    ∃ v, p2 = f v  
  | Output x p =>  
    p2 = p  
  | Sum p1' p2' =>  
    p2 = Sum p1' p2' ∨ p2 = Sum p1 p2'  
  | Res x p =>  
    ∃ y, p2 = Res y p ∧ x ≠ y  
  | _ => False  
end.
```

Inductive **multi_step** : **relation Process** :=

```
| refl : ∀ p, multi_step p p  
| step : ∀ p1 p2 p3, small_step p1 p2 → multi_step p2 p3 → multi_step p1 p3.
```

Inductive **bisim** : **relation Process** :=

```
| bisim_refl : ∀ p, bisim p p  
| bisim_trans : ∀ p1 p2 p3, bisim p1 p2 → small_step p2 p3 → bisim p1 p3  
| bisim_symm : ∀ p1 p2, bisim p1 p2 → bisim p2 p1.
```

Definition bisimulation (*R* : **relation Process**) :=

```
∀ p q, R p q → (∀ p', small_step p p' → ∃ q', R p' q' ∧ small_step q q') ∧
```

$(\forall q', \text{small_step } q \ q' \rightarrow \exists p', R \ p' \ q' \wedge \text{small_step } p \ p')$.

Definition bisimilarity ($R : \text{relation Process}$) :=

$\forall p \ q, R \ p \ q \rightarrow \exists q', \text{bisim } q \ q' \wedge \text{multi_step } p \ q'$.